Original text

The encoding process therefore shall be as follows:

— A group of K information bits $u = [u1, u2, \ldots, uK]$ are collected and copied to the output of the encoder to form a block of systematic code bits. They are also the input to the zero-padding block.

— A total of S zero padding bits are appended at the end of u to form the full-length information bit block $u^* = [u \mid 0, \ldots, 0]$, which is then sent to the information bit de-interleaver module, which in turn produces the bit-de-interleaved sequence $u'' = \pi^{-1}_{info}(u^*)$. π^{-1}_{info} represents the de-interleaver mapping of information bits that permutes u^* to u''.

— The de-interleaved QC-LDPC information bits u'' are sent to the QC-LDPC encoding engine, and used to compute parity-check bits p'' with the parity-check matrix H, and p'' is then interleaved to get $p^* = \pi_{parity}(p'')$. π_{parity} represents the interleaver mapping of parity bits that permutes p'' to p^*.

— $M + P$ parity bits $p^* = [p_1, p_2, \ldots, p_M \mid p_{M+1}, \ldots, p_{M+P}]$ are sent to the puncturing block.

— The last P bits of p^* are truncated, and M parity bits $p = [p_1, p_2, \ldots, p_M]$ are copied to the output of the encoder to form the parity-check bits.

— The FEC codeword without delimiter is $c = [u \mid p] = [u_1, u_2, \ldots, u_K \mid p_1, p_2, \ldots, p_M]$, such that $[u'' \mid p''] H^T = 0$.

Proposed text (changes unmarked)

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— M interleaved parity bits $p^* = [p_1, p_2, \ldots, p_M]$ are copied to the output of the encoder to form the parity-check bits p.

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